

Energy-aware routing in IP networks

Youcef Magnouche, Jérémie Leguay, Feng Zeng

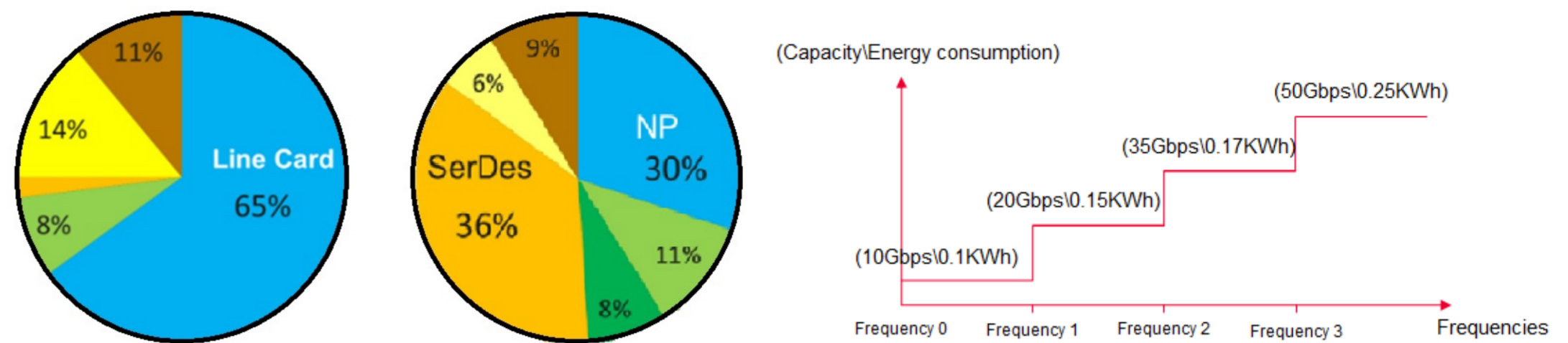


Data Communication Network Algorithm and Measurement Technology Laboratory
Huawei Technologies France, Boulogne-Billancourt, 92100, France

Context & Motivations

- ▶ Information and Communication Technologies represent 10% of the global CO2 emissions
- ▶ 37% of these emissions are due to telecommunication infrastructures and their devices
- ▶ Networking devices are wasting a considerable amount of power
- ▶ Many resources (i.e., routers and links) are powered on without being fully utilized
- ▶ A node (router or switch) is composed of a set of boards. Each board has a set of chipsets on it, every chipset has a set of Serdes (= Port)
- ▶ We propose two methods to optimize the energy consumption

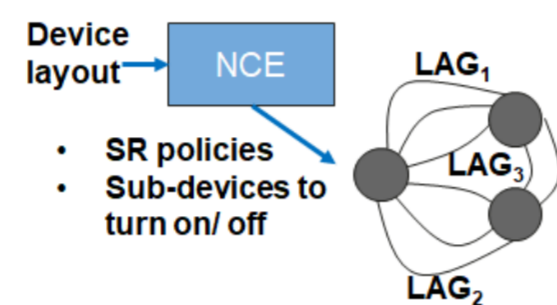
Energy consumption of sub-devices



(Method 1) Turning On/Off Boards, Chipsets & Ports

Select routing paths & sub-devices to turn-off such that:

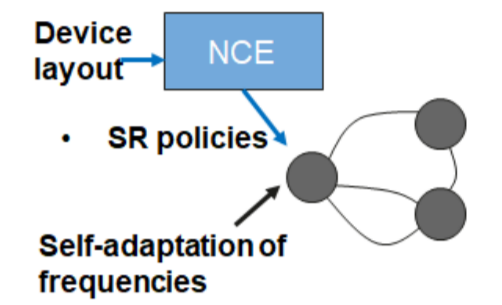
- ▶ Total saved energy is maximum
- ▶ The traffic demand is satisfied
- ▶ Keep at least 1 link in LAGs (avoid IGP disruptions)



(Method 2) Multi-Frequency Chipsets & Serdes

Select routing paths such that:

- ▶ One frequency is selected for every Chipset and Serdes
- ▶ Total saved energy is maximum
- ▶ The traffic demand is satisfied



(Method 1) Mathematical model

$$\begin{aligned} \max \quad & \sum_{d \in D} e_d t_d - \sum_{k \in K} \sum_{p \in P^k} c_p x_p^k && \leftarrow \text{Energy saving} \\ \text{s.t.} \quad & t_d + t_{c_d} + t_{b_d} \leq 1, && \leftarrow \text{Cannot turn-off more than 1 nested sub-device} \\ & \sum_{p \in P^k} x_p^k \geq 1 \quad \forall k \in K, && \leftarrow \text{Demand constraints} \\ & \sum_{k \in K} b(k) \sum_{p \in P^k | c \in p} x_p^k \leq c_e (1 - (t_u + t_{c(u)} + t_{b(u)})), \quad \forall e \in E && \leftarrow \text{Capacity constraints} \\ & \sum_{uv \in T} (t_u + t_{c(u)} + t_{b(u)}) \leq |T| - 1, \quad \forall \text{trunk } T \subseteq E, && \leftarrow \text{Trunk constraint} \\ & t_u + t_{c(u)} + t_{b(u)} = t_v + t_{c(v)} + t_{b(v)}, \quad \forall uv \in E && \leftarrow \text{Link constraints} \\ & t_d \in \{0, 1\}, \quad \forall d \in D, && \leftarrow \text{Integrality constraints} \\ & x_p^k \in \{0, 1\}, \quad \forall k \in K, \forall p \in P^k. && \leftarrow \text{Trivial constraints} \end{aligned}$$

(Method 2) Mathematical model

$$\begin{aligned} \max \quad & \sum_{\text{chipset } d \in D} \sum_{l \in L_d} e_d^l t_d^l - \sum_{k \in K} \sum_{p \in P^k} c_p x_p^k && \leftarrow \text{Energy saving} \\ \text{s.t.} \quad & \sum_{p \in P^k} x_p^k \geq 1 \quad \forall k \in K, && \leftarrow \text{Demand constraints} \\ & \sum_{l \in L_d} y_d^l = 1 \quad \forall d \in D, && \leftarrow \text{Frequency selection constraint} \\ & \sum_{k \in K} b(k) \sum_{p \in P^k | T \cap p = \emptyset} x_p^k \leq \sum_{l \in L_d (p_1, p_2) \in T} c_{p_1}^l t_{p_1}^l, \quad \forall \text{trunk } T \subseteq E && \leftarrow \text{Serdes Capacity constraint} \\ & \sum_{l \in L_d} b(k) \sum_{\text{serdes } p \in d} c_p^l t_p^l \leq \sum_{l \in L_d} c_d^l t_d^l, \quad \forall \text{chipset } d \subseteq D && \leftarrow \text{Chipset Capacity constraint} \\ & t_d^l \in \{0, 1\}, \quad \forall d \in D, l \in L_d, && \leftarrow \text{Integrality constraints} \\ & x_p^k \in \{0, 1\}, \quad \forall k \in K, \forall p \in P^k. && \leftarrow \text{Trivial constraints} \end{aligned}$$

(Method 1) Advantages

- ▶ Compatible with current routers
- ▶ More energy saving thanks to the shut-down of sub-devices (even boards)

(Method 2) Advantages

- ▶ Fast adaptation to traffic in all scenarios
- ▶ Local energy saving mechanism (not controlled by NCE)

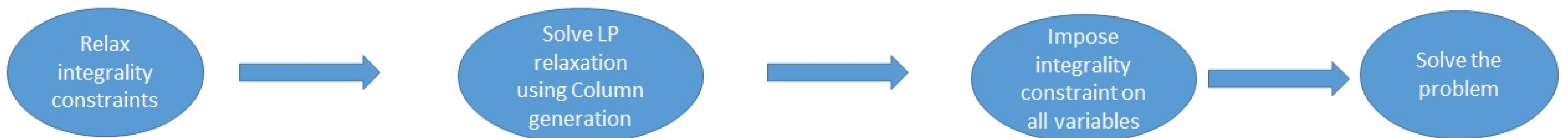
(Method 1) Drawbacks

- ▶ Needs accurate traffic predictions
- ▶ Works mainly for networks with lots of LAGs
- ▶ Large reaction time in case of anomalies (wake-up time)

(Method 2) Drawbacks

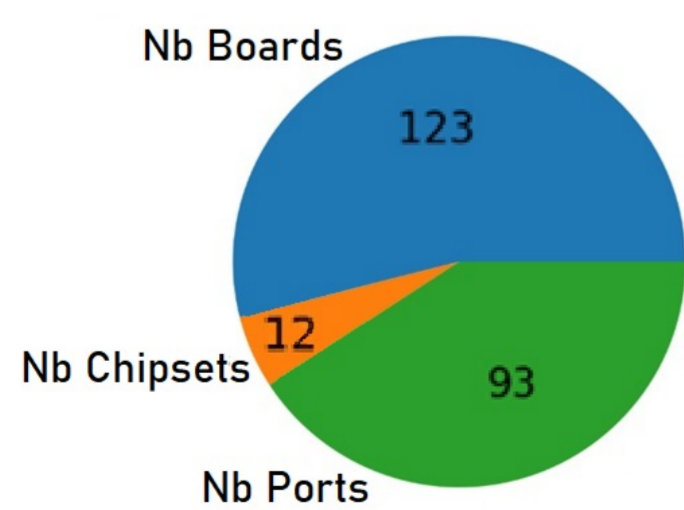
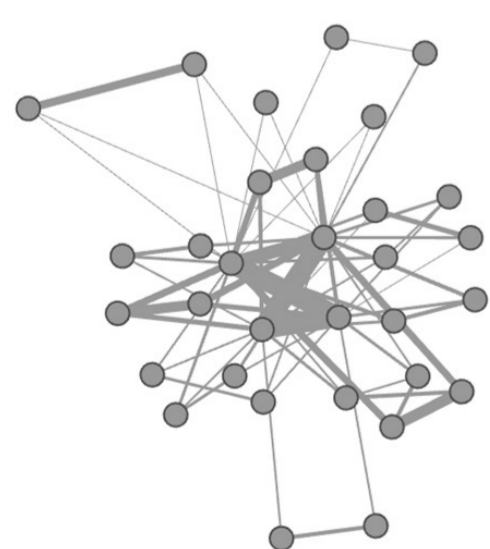
- ▶ Sub-devices cannot be shut-down totally
- ▶ Needs frequent accurate traffic measurements
- ▶ Compatible only with last generation Chipsets and Serdes (does not work with boards)

Column generation based heuristic



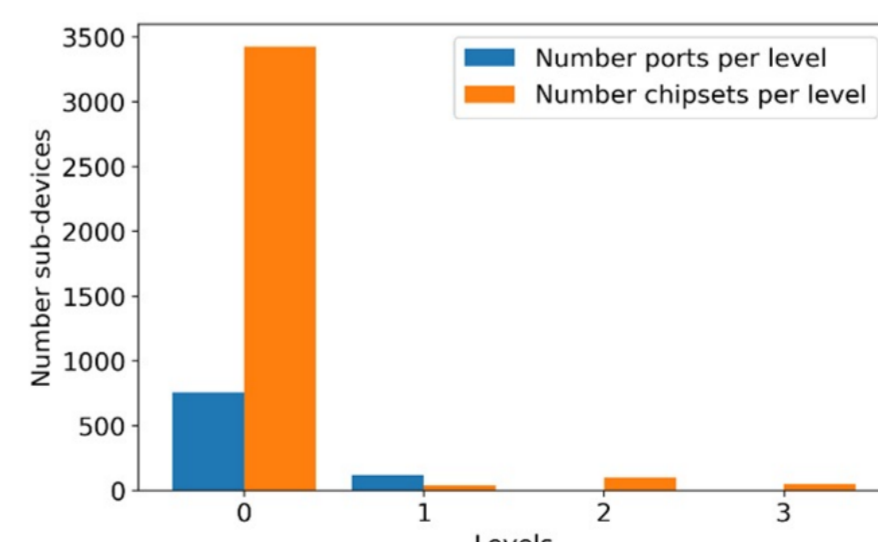
Telecom operator A (32 nodes, 992 demands, 40% of maximum traffic)

Turning On/Off Boards, Chipsets & Ports



- ▶ Initial energy consumption: 323500.0 W
- ▶ Total saved energy: 55210.4 W
- ▶ Energy saved Gap: 17.07%
- ▶ Number of turned-off links: 274
- ▶ France: 48.36 Tons CO2/y
- ▶ Europe: 217.64 Tons CO2/y
- ▶ China: 464.30 Tons CO2/y

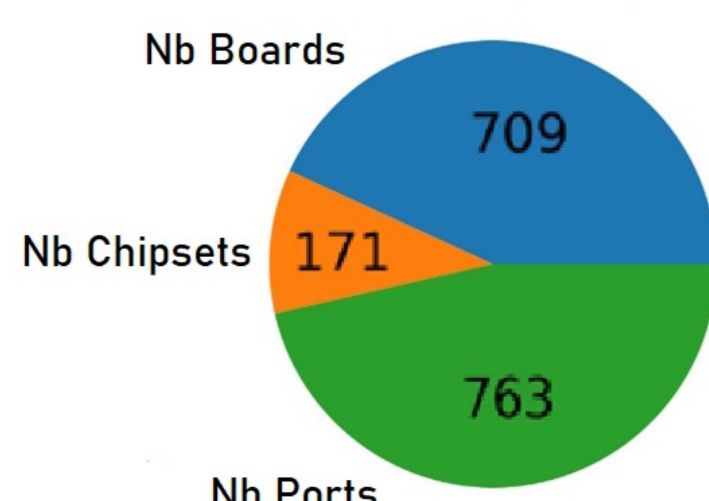
Multi-Frequency Chipsets & Serdes



- ▶ Initial energy consumption: 323500.0 W
- ▶ Total saved energy: 37057.2 W
- ▶ Energy saved Gap: 11.46%
- ▶ Number of turned-off links: 0
- ▶ France: 32.46 Tons CO2/y
- ▶ Europe: 146.08 Tons CO2/y
- ▶ China: 311.64 Tons CO2/y

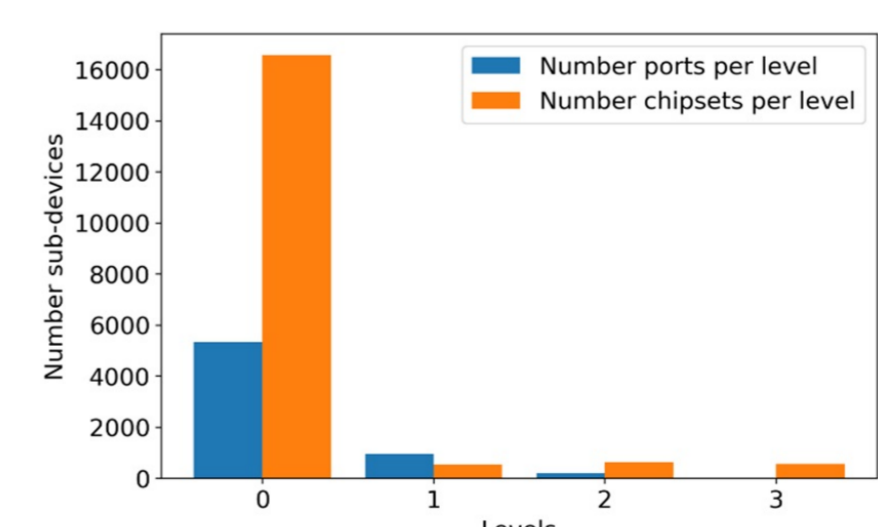
Telecom operator B (143 nodes, 20306 demands, 40% of maximum traffic)

Turning On/Off Boards, Chipsets & Ports



- ▶ Initial energy consumption: 2408870 W
- ▶ Total saved energy: 361128.0 W
- ▶ Energy saved Gap: 14.99%
- ▶ Number of turned-off links: 2292
- ▶ France: 316.35 Tons CO2/y
- ▶ Europe: 1423.57 Tons CO2/y
- ▶ China: 3036.94 Tons CO2/y

Multi-Frequency Chipsets & Serdes



- ▶ Initial energy consumption: 2408870 W
- ▶ Total saved energy: 300851.0 W
- ▶ Energy saved Gap: 12.49%
- ▶ Number of turned-off links: 0
- ▶ France: 263.55 Tons CO2/y
- ▶ Europe: 1185.95 Tons CO2/y
- ▶ China: 2530.04 Tons CO2/y